

Research Article

# Complexity of Some Types of Cyclic Snake Graphs

**Basma Mohamed\***, **Mohamed Amin**

Mathematics and Computer Science Department, Faculty of Science, Menoufia University, Shebin Elkom, Egypt

## Abstract

The number of spanning trees in graphs (networks) is a crucial invariant, and it is also an important measure of the reliability of a network. Spanning trees are special subgraphs of a graph that have several important properties. First,  $T$  must span  $G$ , which means it must contain every vertex in graph  $G$ , if  $T$  is a spanning tree of graph  $G$ .  $T$  needs to be a subgraph of  $G$ , second. Stated differently, any edge present in  $T$  needs to be present in  $G$  as well. Third,  $G$  is the same as  $T$  if each edge in  $T$  is likewise present in  $G$ . In path-finding algorithms like Dijkstra's shortest path algorithm and  $A^*$  search algorithm, spanning trees play an essential part. In those approaches, spanning trees are computed as component components. Protocols for network routing also take advantage of it. In numerous techniques and applications, minimum spanning trees are highly beneficial. Computer networks, electrical grids, and water networks all frequently use them. They are also utilized in significant algorithms like the min-cut max-flow algorithm and in graph issues like the travelling salesperson problem. In this paper, we use matrix analysis and linear algebra techniques to obtain simple formulas for the number of spanning trees of certain kinds of cyclic snake graphs.

## Keywords

Cyclic Snakes, Subdivision, Edge Contraction, Spanning Trees

## 1. Introduction

In this paper, we give some basic definitions. We deal with simple and finite undirected graphs  $G = (V, E)$ , where  $V$  is the vertex set and  $E$  is the edge set. For a graph  $G$ , a spanning tree in  $G$  is a tree that has the same vertex set as  $G$ . The number of spanning trees in  $G$ , also called the complexity of the graph, denoted by  $\tau(G)$ , is a well-studied quantity (for a long time) and appears in a number of applications. The most notable application fields are network reliability [1–3], enumerating certain chemical isomers [4], and counting the number of Eulerian circuits in a graph [5]. A classical result of Kirchhoff [6] can be used to determine the number of spanning trees for  $G = (V, E)$ . Let  $V = \{v_1, v_2, \dots, v_n\}$ ,

then the Kirchhoff matrix  $H$  defined as  $n \times n$  characteristic matrix  $H = D - A$ , where  $D$  is the diagonal matrix of the degrees of  $G$  and  $A$  is the adjacency matrix of  $G$ ,  $H = [a_{ij}]$  defined as follows: (i)  $a_{ij} = -1$ ,  $v_i$  and  $v_j$  are adjacent and  $i \neq j$ , (ii)  $a_{ij}$  equals the degree of vertex  $v_i$  if  $i = j$ , and (iii)  $a_{ij} = 0$  otherwise. All of the co-factors of  $H$  are equal to  $\tau(G)$ . There are other methods for calculating  $\tau(G)$ . Let  $\mu_1 \geq \mu_2 \geq \dots \geq \mu_p$  denote the eigenvalues of  $H$  matrix of  $p$  point graph. Then it is easily shown that  $\mu_p = 0$ . Furthermore, Kelmans and Chelnokov [7] have shown that,  $\tau(G) = 1/p \prod_{k=1}^{p-1} \mu_k$ . The formula for the number of span-

\*Corresponding author: [bosbos25jan@yahoo.com](mailto:bosbos25jan@yahoo.com) (Basma Mohamed)

**Received:** 5 December 2023; **Accepted:** 26 December 2023; **Published:** 20 February 2024



ning trees in a  $d$ -regular graph  $G$  can be expressed as  $\tau(G) = 1/p \prod_{k=1}^{p-1} (d - \mu_k)$  where  $\lambda_0 = d, \lambda_1, \lambda_2, \dots, \lambda_{p-1}$  are the eigenvalues of the corresponding adjacency matrix of the graph. However, for a few special families of graphs, there exist simple formulas that make it much easier to calculate and determine the number of corresponding spanning trees, especially when these numbers are very large.

One of the first such results is due to Cayley [8], who showed that for a complete graph with  $n$  vertices,  $K_n$  has  $n^{n-2}$  spanning trees that he showed  $\tau(K_n) = n^{n-2}, n \geq 2$ .

$\tau(K_{p,q}) = p^{q-1} q^{p-1}$   $p, q \geq 1$  where  $K_{p,q}$  is the complete bipartite graph with bipartite sets containing  $p$  and  $q$  vertices, respectively. It is well known, as in e.g., [9, 10]. Another result is due to Sedlacek [11], who derived a formula for the wheel on  $n + 1$  vertices,  $W_{n+1}$ , which is formed from a cycle  $C_n$  on  $n$  vertices by adding a vertex adjacent to every vertex of  $C_n$ . In particular, he showed that  $\tau(W_{n+1}) = (3 + \sqrt{5}/2)^n + (3 - \sqrt{5}/2)^n - 2$ , for  $n \geq 3$ . Sedlacek [12] also derived a formula for the number of spanning trees in a Mobius ladder. The Mobius ladder  $M_n$  is formed from cycle  $C_{2n}$  on  $2n$  vertices labelled  $v_1, v_2, \dots, v_{2n}$  by adding an edge  $v_i v_{i+n}$  for every vertex  $v_i$  where  $i \leq n$ .

The number of spanning trees in  $M_n$  is given by  $\tau(M_n) = n/2[(2 + \sqrt{3})^n + (2 - \sqrt{3})^n + 2]$  for  $n \geq 2$ . Another class of graphs for which an explicit formula has been derived is based on a prism [13, 14]. Let the vertices of two disjoint length cycles be labelled  $v_1, v_1, \dots, v_n$  in one cycle and  $w_1, w_1, \dots, w_n$  in the other. The prism  $R_n$  is defined as the graph obtained by adding to these two cycles all the edges of the form  $v_i w_i$ . The number of spanning trees in  $R_n$  is given by the following formula  $n/2[(2 + \sqrt{3})^n + (2 - \sqrt{3})^n - 2]$ . For more results, it is suggested to see these articles [15-28].

Chio's condensation is a method for evaluating an  $n \times n$  determinant in terms of  $(n - 1) \times (n - 1)$  determinants; see [29]:

$$A = \begin{vmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{vmatrix} = \begin{vmatrix} \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} & \begin{vmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{vmatrix} & \dots & \begin{vmatrix} a_{11} & a_{1n} \\ a_{21} & a_{2n} \end{vmatrix} \\ \begin{vmatrix} a_{11} & a_{12} \\ a_{31} & a_{32} \end{vmatrix} & \begin{vmatrix} a_{11} & a_{13} \\ a_{31} & a_{33} \end{vmatrix} & \dots & \begin{vmatrix} a_{11} & a_{1n} \\ a_{31} & a_{3n} \end{vmatrix} \\ \vdots & \vdots & \ddots & \vdots \\ \begin{vmatrix} a_{11} & a_{12} \\ a_{n1} & a_{n2} \end{vmatrix} & \begin{vmatrix} a_{11} & a_{13} \\ a_{n1} & a_{n3} \end{vmatrix} & \dots & \begin{vmatrix} a_{11} & a_{1n} \\ a_{n1} & a_{nn} \end{vmatrix} \end{vmatrix}$$

## 2. Dodgson's Condensation Method

Dodgson's condensation method computes determinants

of size  $n \times n$  by expressing them in terms of those of size  $(n - 1) \times (n - 1)$ , and then expresses the latter in terms of determinants of size  $(n - 2) \times (n - 2)$ , and so on (see [30]).

This method is based on Dodgson and Chio's method, but the difference between them is that this new method is resolved by calculating 4 unique determinants of  $(n - 1) \times (n - 1)$  order, (which can be derived from determinants of  $n \times n$  order, if we remove first row and first column or first row and last column or last row and first column or last row and last column, elements that belong to only one of the unique determinants), and one determinant of  $(n - 2) \times (n - 2)$  order which is formed from  $n \times n$  order determinant with elements  $a_{ij}$  with  $i, j = 1, n$ , on condition that the determinant of  $(n - 2) \times (n - 2) = 0$ .

$$A = \frac{1}{\begin{vmatrix} a_{22} & \dots & a_{2,n-1} \\ \vdots & \ddots & \vdots \\ a_{n-1,2} & \dots & a_{n-1,n-1} \end{vmatrix}} \begin{pmatrix} \begin{vmatrix} a_{11} & \dots & a_{1,n-1} \\ \vdots & \ddots & \vdots \\ a_{n-1,1} & \dots & a_{n-1,n-1} \end{vmatrix} & \begin{vmatrix} a_{12} & \dots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{n-1,2} & \dots & a_{n-1,n} \end{vmatrix} \\ \vdots & \vdots \\ \begin{vmatrix} a_{21} & \dots & a_{2,n-1} \\ \vdots & \ddots & \vdots \\ a_{n,1} & \dots & a_{n,n-1} \end{vmatrix} & \begin{vmatrix} a_{22} & \dots & a_{2n} \\ \vdots & \ddots & \vdots \\ a_{n,2} & \dots & a_{nn} \end{vmatrix} \end{pmatrix}$$

## 3. Complexity of Some Types of Cyclic Snake Graphs

Lemma 4.1: Given a simple graph  $G$  with  $n$  vertices, its laplacian matrix  $L_{n \times n}$  is defined as:  $L = D - A$  where  $D$  is the degree matrix and  $A$  is the adjacency matrix of the graph. In the case of directed graphs, either the in degree or the out degree might be used, depending on the application. The elements of  $L$  are given by

$$L_{i,j} := \begin{cases} \deg(v_i) & \text{if } i = j \\ -1 & \text{if } i \neq j \text{ and } v_i \text{ is adjacent to } v_j \\ 0 & \text{otherwise} \end{cases}$$

where  $\deg(v_i)$  is degree of the vertex  $i$

Definition 4.1 [31]:

A triangular snake (or  $\Delta_k$ -snake) is a connected graph in which all blocks are triangles and the block-cut-point graph is a path.

Definition 4.2: The  $kc_n$ -snake is called linear, if the block-cut-vertex graph of  $kc_n$ -snake has the property that the distance between any two consecutive cut-vertices is

$\lfloor \frac{n}{2} \rfloor$  Definition 4.3: The graphs  $(m, k) C_4$  as the family of

graphs  $kc_4$ -snake where every block has  $m$  copies of  $C_4$  with two non- adjacent vertices in common, where the number of blocks is denoted by  $k$ .

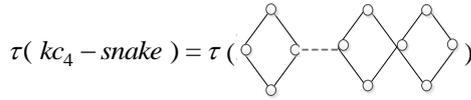
Definition 4.4: The subdivision of a graph  $G$  is obtained by subdividing every edge of  $G$  exactly once.

The Main Results:

Theorem 5.1 The number of spanning trees of the linear

$kc_4 - snake$

$$\tau(G) = \frac{1}{2} \begin{bmatrix} 3 & 1 \\ 1 & 3 \end{bmatrix} = 4.$$



$$\tau(kc_4 - snake) = \tau(\text{graph})$$

We try to prove at  $m=k$  that is Straightforward induction using properties of determinants and Dodgson and Chio method and applying Lemma 2.1, we

$\tau(kc_4 - snake) = 4^k, k \geq 2$ , where  $k$  is the number of blocks. have:

Proof. By induction, prove at  $m = 1$ ,  $\tau(G)$  is a  $2 \times 2$  matrix with both rows the same:

$$\tau(G) = \begin{bmatrix} 4 & 0 & \dots & \dots & \dots & \dots & \dots & 0 & -1 & -1 & -1 & 0 \\ 0 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & -1 & -1 & -1 & -1 & 0 \\ \vdots & \ddots & 4 & -1 & \ddots & -1 & -1 & -1 & 0 & 0 & \vdots & \vdots \\ \vdots & \ddots & -1 & 2 & -1 & \ddots & 0 & 0 & 0 & \vdots & \vdots & \vdots \\ \vdots & \ddots & 0 & -1 & \ddots & -1 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots \\ \vdots & \ddots & -1 & 0 & -1 & \ddots & 0 & \ddots & \ddots & \ddots & \vdots & \vdots \\ \vdots & -1 & -1 & 0 & \dots & 0 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots \\ 0 & -1 & -1 & 0 & \dots & \dots & \dots & \ddots & \ddots & \ddots & \ddots & \ddots \\ -1 & -1 & 0 & \dots & \dots & \dots & \dots & \ddots & \ddots & \ddots & \ddots & \ddots \\ -1 & -1 & \vdots & \dots & \dots & \dots & \dots & \ddots & \ddots & \ddots & \ddots & 0 \\ -1 & 0 & \vdots & \dots & \dots & \dots & \dots & \ddots & \ddots & \ddots & \ddots & -1 \\ 0 & \vdots & 0 & \dots & -1 & 2 \end{bmatrix}_{2k \times 2k}$$

$$= \frac{1}{2 * 4^{k-1}} \begin{vmatrix} 3 * 4^{k-1} & -4^{k-1} \\ -4^{k-1} & 3 * 4^{k-1} \end{vmatrix} = \frac{(3 * 4^{k-1})^2 - (-4^{k-1})^2}{2 * 4^{k-1}} = \frac{8 * 4^{2k-2}}{2 * 4^{k-1}} = 4^k, k \geq 2, \text{Where } k \text{ is the number of blocks.}$$

$$S = \begin{bmatrix} 4 & 0 & \dots & \dots & \dots & \dots & \dots & 0 & -1 & -1 & -1 & -1 & 0 \\ 0 & \ddots & \ddots & \ddots & \dots & \dots & \dots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & 4 & -1 & \ddots & -1 & -1 & -1 & 0 & \dots & \dots & \dots & \vdots \\ \vdots & \ddots & -1 & 2 & -1 & \ddots & 0 & \dots & \dots & \dots & \dots & \dots & \vdots \\ \vdots & \ddots & \ddots & -1 & \ddots & -1 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \dots & -1 & \ddots & -1 & \ddots & 0 & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & -1 & 0 & \ddots & 0 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & \ddots & -1 & \vdots & \ddots & \vdots \\ -1 & \ddots & 0 & \vdots & \ddots & \vdots \\ -1 & \ddots & \vdots & \ddots & \vdots \\ -1 & \ddots & \vdots & \ddots & \vdots \\ -1 & \ddots & \vdots & \ddots & 0 \\ 0 & \dots & 0 & 2 \end{bmatrix}_{(2k-2) \times (2k-2)}$$

$$= 2 * 4^{k-1}$$

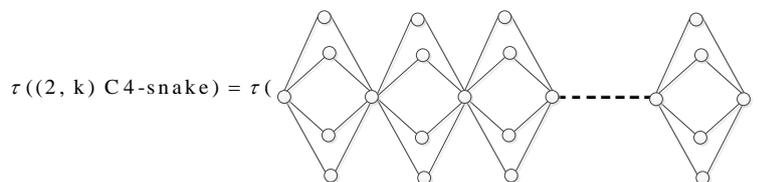
$$A = \begin{bmatrix} 4 & 0 & \dots & 0 & -1 & -1 & -1 \\ 0 & \ddots & \ddots & \dots & \dots & \dots & \dots & \dots & \dots & -1 & -1 & -1 & 0 \\ \vdots & \ddots & 4 & -1 & 0 & -1 & -1 & -1 & 0 & \dots & \dots & \dots & \vdots \\ \vdots & \vdots & -1 & 2 & -1 & \ddots & \dots & \dots & \dots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \ddots & -1 & \ddots & -1 & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & -1 & \ddots & -1 & \ddots & 0 & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & -1 & \vdots & \ddots & 0 & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & -1 & \vdots & \dots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & 0 & \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots \\ \vdots & -1 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & -1 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ -1 & -1 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ -1 & -1 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & 0 \\ -1 & 0 & \dots & 0 & 2 \end{bmatrix}_{(2k-2) \times (2k-2)}$$

$$= 3 * 4^{k-1},$$

$$B = B^T = \begin{pmatrix} 0 & \dots & 0 & -1 & -1 & -1 & 0 \\ 4 & \ddots & \dots & \dots & \dots & \dots & \dots & \dots & 0 & -1 & -1 & -1 & -1 & 0 & \vdots \\ 0 & \ddots & \ddots & \dots & \dots & 0 & -1 & -1 & -1 & -1 & 0 & \dots & \dots & \vdots & \vdots \\ \vdots & \ddots & 4 & -1 & 0 & -1 & -1 & -1 & 0 & \dots & \dots & \dots & \dots & \vdots & \vdots \\ \vdots & \vdots & -1 & 2 & -1 & 0 & \dots & \vdots & \vdots \\ \vdots & \vdots & 0 & -1 & \ddots & -1 & \ddots & \dots & \dots & \dots & \dots & \dots & \dots & \vdots & \vdots \\ \vdots & \vdots & -1 & 0 & -1 & \ddots & 0 & \dots & \dots & \dots & \dots & \dots & \dots & \vdots & \vdots \\ \vdots & \vdots & -1 & \vdots & 0 & 0 & \ddots & \vdots & \vdots \\ 0 & -1 & -1 & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ -1 & -1 & 0 & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ -1 & -1 & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ -1 & -1 & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ -1 & 0 & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & \vdots & \vdots & \vdots & \vdots & \dots & \dots & \dots & \dots & \dots & \dots & 0 & 2 & -1 & \dots \end{pmatrix}_{(2k-2) \times (2k-2)} = 4^{k-1}$$

$$C = \begin{pmatrix} 4 & 0 & \dots & \dots & \dots & \dots & \dots & 0 & -1 & -1 & -1 & -1 & 0 & 0 \\ 0 & \ddots & \ddots & \ddots & \ddots & \ddots & -1 & -1 & -1 & -1 & 0 & \dots & \dots & \vdots \\ \vdots & \ddots & 4 & -1 & 0 & -1 & -1 & -1 & 0 & \dots & \dots & \dots & \dots & \vdots \\ \vdots & \ddots & -1 & 2 & -1 & \ddots & \dots & \vdots \\ \vdots & \ddots & 0 & -1 & \ddots & -1 & \ddots & \vdots \\ \vdots & \ddots & -1 & \ddots & -1 & \ddots & 0 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & -1 & -1 & \vdots & \ddots & 0 & \ddots & \vdots \\ 0 & -1 & -1 & \vdots & \vdots & \vdots & \ddots & \vdots \\ -1 & -1 & 0 & \vdots & \vdots & \vdots & \ddots & \vdots \\ -1 & -1 & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ -1 & 0 & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ -1 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & -1 \\ 0 & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & 0 & -1 & 2 \end{pmatrix}_{(2k-2) \times (2k-2)} = 3 \cdot 4^{k-1}$$

Theorem 5.2 The number of spanning trees of the  $(2, k)$  C4-snake



$$\tau((2, k) \text{ C4-snake}) = (32)^k k \geq 3$$

Proof. By induction, prove at  $m = 1$ ,  $\tau(G)$  is a  $2 \times 2$  matrix with both rows the same:

$$\tau(G) = \frac{1}{24} \begin{bmatrix} 32 & 16 \\ 16 & 32 \end{bmatrix} = 32.$$

We try to prove at  $m=k$  that is

Straightforward induction using properties of determinants and Dodgson and Chio method and applying Lemma 2.1, we have:









## Conflicts of Interest

The authors declare no conflicts of interest.

## References

- [1] Colbourn, C. J. (1987). *The combinatorics of network reliability*. Oxford University Press, Inc.
- [2] Myrvold, W., Cheung, K. H., Page, L. B., & Perry, J. E. (1991). Uniformly-most reliable networks do not always exist. *Networks*, 21(4), 417-419.
- [3] Petingi, L., Boesch, F., & Suffel, C. (1998). On the characterization of graphs with maximum number of spanning trees. *Discrete mathematics*, 179(1-3), 155-166.
- [4] Brown, T. J., Mallion, R. B., Pollak, P., & Roth, A. (1996). Some methods for counting the spanning trees in labelled molecular graphs, examined in relation to certain fullerenes. *Discrete Applied Mathematics*, 67(1-3), 51-66.
- [5] Moon, J. W. (1967). Enumerating labelled trees. *Graph theory and theoretical physics*, 261271.
- [6] Kirchhoff, G. (1847). Ueber die Auflösung der Gleichungen, auf welche man bei der Untersuchung der linearen Vertheilung galvanischer Ströme geführt wird. *Annalen der Physik*, 148(12), 497-508.
- [7] Kelmans, A. K., & Chelnokov, V. M. (1974). A certain polynomial of a graph and graphs with an extremal number of trees. *Journal of Combinatorial Theory, Series B*, 16(3), 197-214.
- [8] Cayley, G. A. (1889). A Theorem on trees, *Quarterly Journal of Mathematics*, 23, 276-378.
- [9] Clark, L. (2003). On the enumeration of multipartite spanning trees of the complete graph. *Bull. of the ICA*, 38, 50-60.
- [10] Qiao, S. N., & Chen, B. (2009). The number of spanning trees and chains of graphs. *Applied Mathematics E-Notes*, 9, 10-16.
- [11] Sedlacek, J. (1970). On the skeletons of a graph or digraph. In *Proc. Calgary International Conference on Combinatorial Structures and their Applications*, Gordon and Breach (pp. 387-391).
- [12] Sedláček, J. (1970). Lucas number in graph theory. *Mathematics (Geometry and Graph theory) (Czech)*, Univ. Karlova, Prague, 111-115.
- [13] Boesch, F. T., & Bogdanowicz, Z. R. (1987). The number of spanning trees in a prism. *International journal of computer mathematics*, 21(3-4), 229-243.
- [14] Boesch, F. T., & Prodinger, H. (1986). Spanning tree formulas and Chebyshev polynomials. *Graphs and Combinatorics*, 2(1), 191-200.
- [15] S. N. Daoud, The deletion-contraction method for counting the number of spanning trees of graphs, *Eur. Phys. J. Plus*, 130, 2015, 1-14.
- [16] S. N. Daoud, Complexity of graphs generated by wheel graph and their asymptotic limits, *JOEMS*, 25(4), 2017, 424-433.
- [17] S. N. Daoud, Number of spanning trees in different products of complete and complete tripartite graphs, *Ars Comb.*, 139, 2018, 85-103.
- [18] S. N. Daoud, Number of Spanning trees of cartesian and composition products of graphs and Chebyshev polynomials, *IEEE Access*, 7, 2019, 71142-71157.
- [19] J. B. Liu, and S. N. Daoud, The complexity of some classes of pyramid graphs created from a gear graph, *Symmetry*, 10(12), 2018.
- [20] S. N. Daoud and W. Saleh, Number of Spanning Trees of Some of Pyramid Graphs Generated by a Wheel Graph, *Math. Comb*, 43, 2020.
- [21] S. N. Daoud, Complexity of join and corona graphs and Chebyshev polynomials, *Journal of Taibah University for Science*, 12(5), 2018, 557-572.
- [22] B. Mohamed, L. Mohaisen and M. Amin, Binary Equilibrium Optimization Algorithm for Computing Connected Domination Metric Dimension Problem, *Sci. Program*, 2022.
- [23] B. Mohamed and M. Amin, The Metric Dimension of Subdivisions of Lilly Graph, Tadpole Graph and Special Trees," *Appl. Comput. Math*, 12(1), 2023, 9-14.
- [24] B. Mohamed, L. Mohaisen and M. Amin, Computing Connected Resolvability of Graphs Using Binary Enhanced Harris Hawks Optimization, *Intell. Autom. Soft Comput*, 36(2), 2023.
- [25] B. Mohamed, Metric Dimension of Graphs and its Application to Robotic Navigation, *IJCA*, 184(15), 2022.
- [26] B. Mohamed and M. Amin, Domination Number and Secure Resolving Sets in Cyclic Networks, *Appl. Comput. Math*, 12(2), 2023, 42-45.
- [27] B. Mohamed and M. Amin, A hybrid optimization algorithms for solving metric dimension problem, (GRAPH-HOC), 15(1), 2023, 1-10.
- [28] B. Mohamed, A Comprehensive Survey on the Metric Dimension Problem of Graphs and Its Types, *Int. J. Theor. Appl. Math*, 9(1), 2023, 1-5.
- [29] Chi ó F. (1853). *Mémoire sur les fonctions connues sous le nom De Résultantes Ou De Déterminans*. áriteur inconnu.
- [30] Gjonbalaj, Q., & Salihu, A. (2010). Computing the determinants by reducing the orders by four. *Applied Mathematics E-Notes*, 10, 151-158.
- [31] Gallian, J. A. "Dynamic Survey DS6: Graph Labeling." *Electronic J. Combinatorics*, DS6, 1-58, Jan. 3, 2007.